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Unique column combinations

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Guest lecture in Data Profiling and Data Cleansing
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Agenda

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Introduction and problem statement

- Unique column combinations
- Exponential search space
- Null values
- General pruning techniques

Discovery algorithms

- Apriori
- HCA
- DUCC
- Gordian

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Unique column combinations

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- Relational model
 - Dataset R with schema S
- Unique column combination $K \subseteq S$
$$\forall r_i, r_j \in R : i \neq j \Rightarrow r_i[K] \neq r_j[K]$$
- In the following, they are called uniques
- Examples: all primary keys, all unique constraints

A	B	C
a	1	x
b	2	x
c	2	y

- Uniques: { A, AB, AC, BC, ABC }
- Non-uniques: { B, C }

Minimal uniques

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- We are mostly interested in minimal uniques $K \subseteq S$
 $\neg \exists K' \subseteq S : \text{unique}(K') \wedge K' \subset K$
- Removal of any column leads to non-unique combination
- For the previous example: {A, BC}
- Redundant: {AB, AC, ABC}

A	B	C
a	1	x
b	2	x
c	2	y

- Candidates for primary keys

Maximal non-uniques

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- Analogously we can define maximal non-uniques $K \subseteq S$
 $\neg \exists K' \subseteq S : \text{non-unique}(K') \wedge K \subset K'$
- Adding any column leads to unique combination
- Non-unique: {AB, AC}
- Redundant: {A, B, C}

A	B	C
a	1	x
a	2	x
a	2	y

- May be a data quality problem

Applications

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- Learning characteristics about a new data set
- Database management
 - Finding a primary key
 - Finding unique constraints
- Query optimization
 - Cardinality estimations for joins
- Finding duplicates / data quality issues
 - If expected unique column combinations are not unique
 - Or with approximate uniques

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Introduction and problem statement

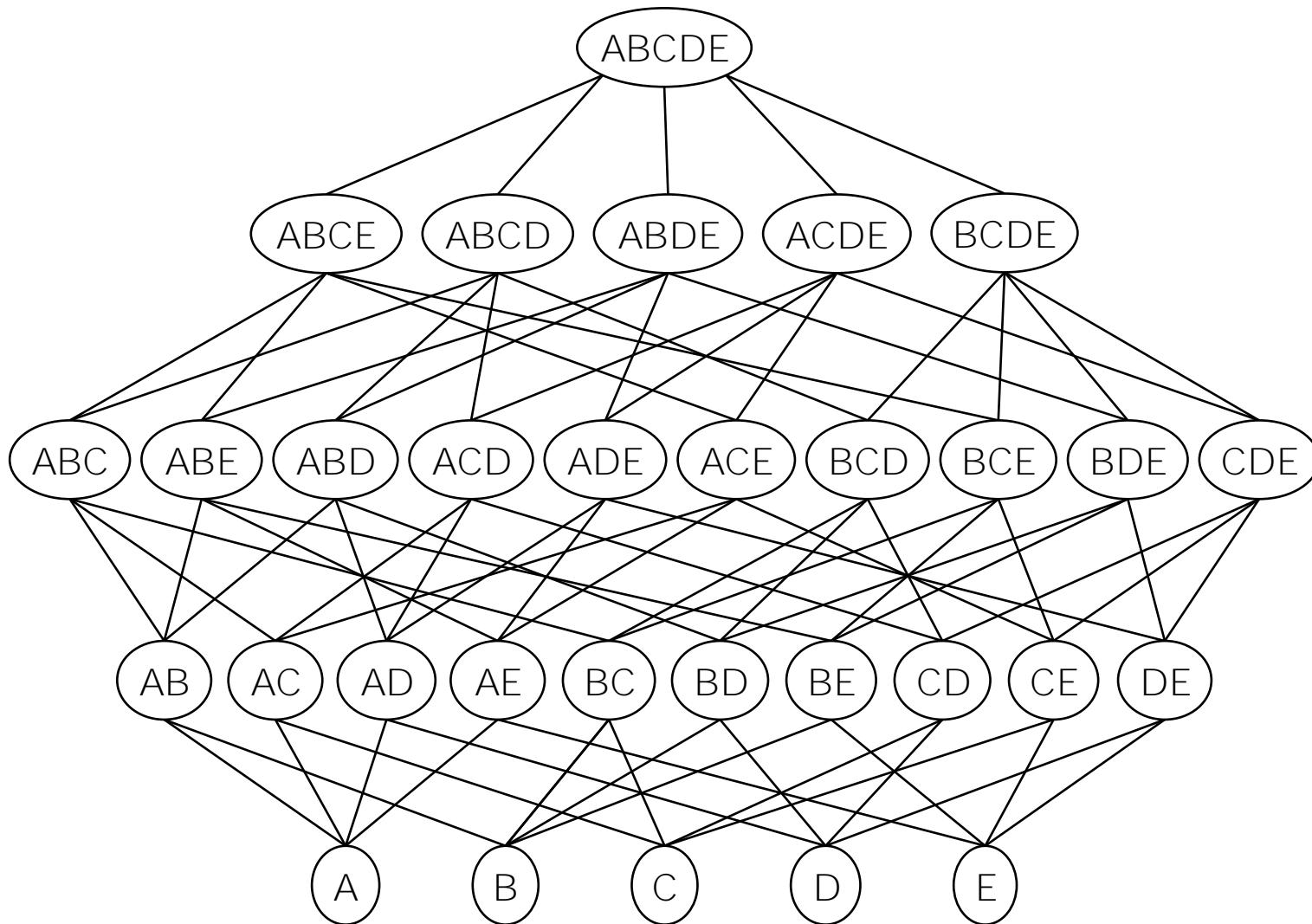
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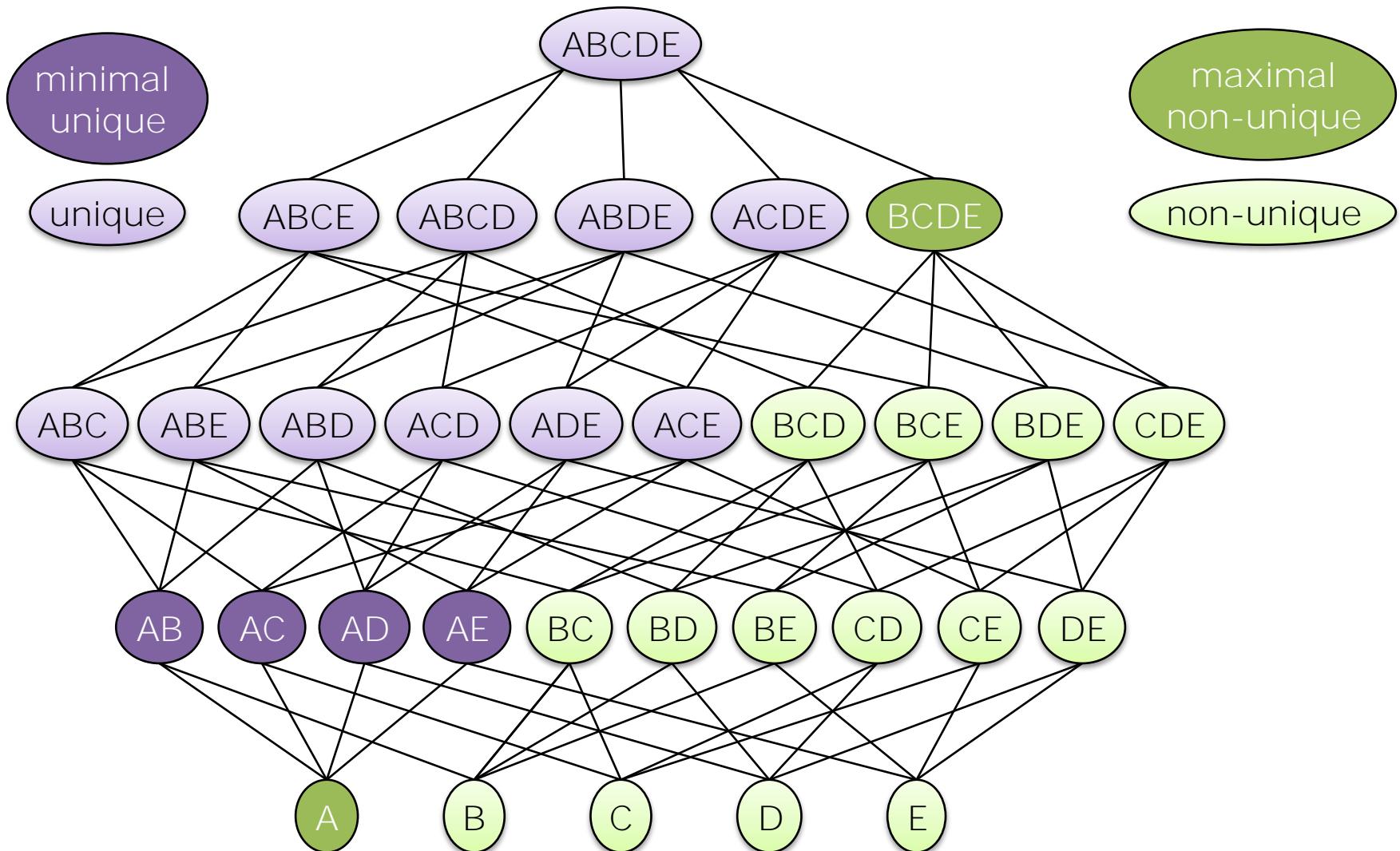
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Exponential search space

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Result of algorithm



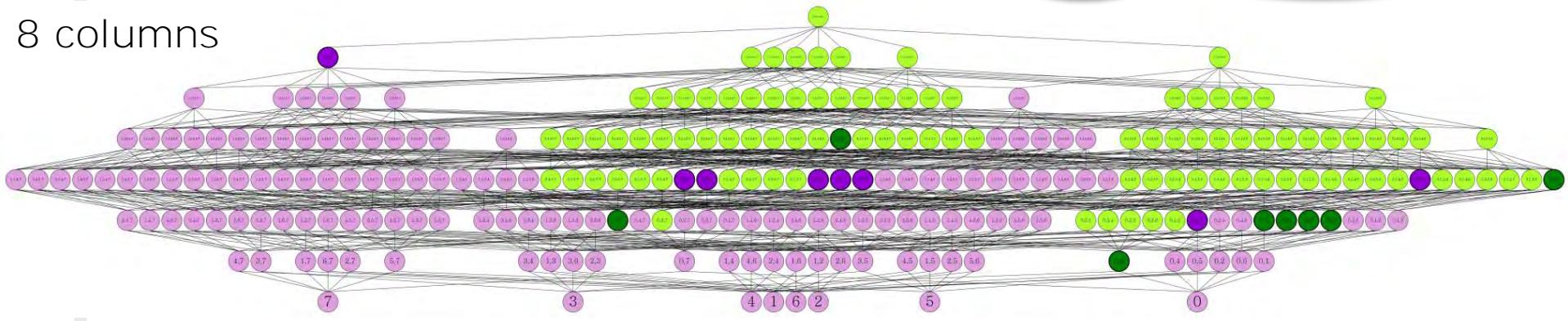
TPCH line item

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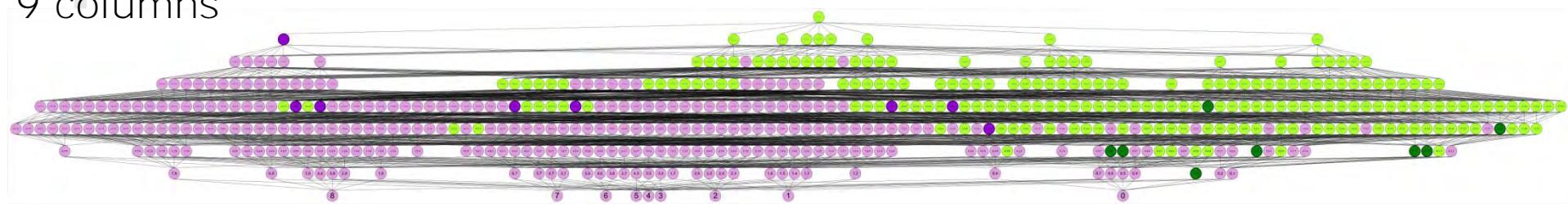
8 columns

unique

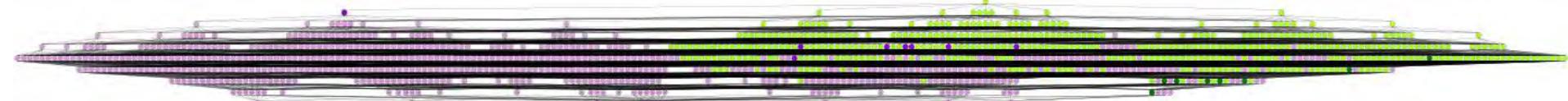
non-unique



9 columns

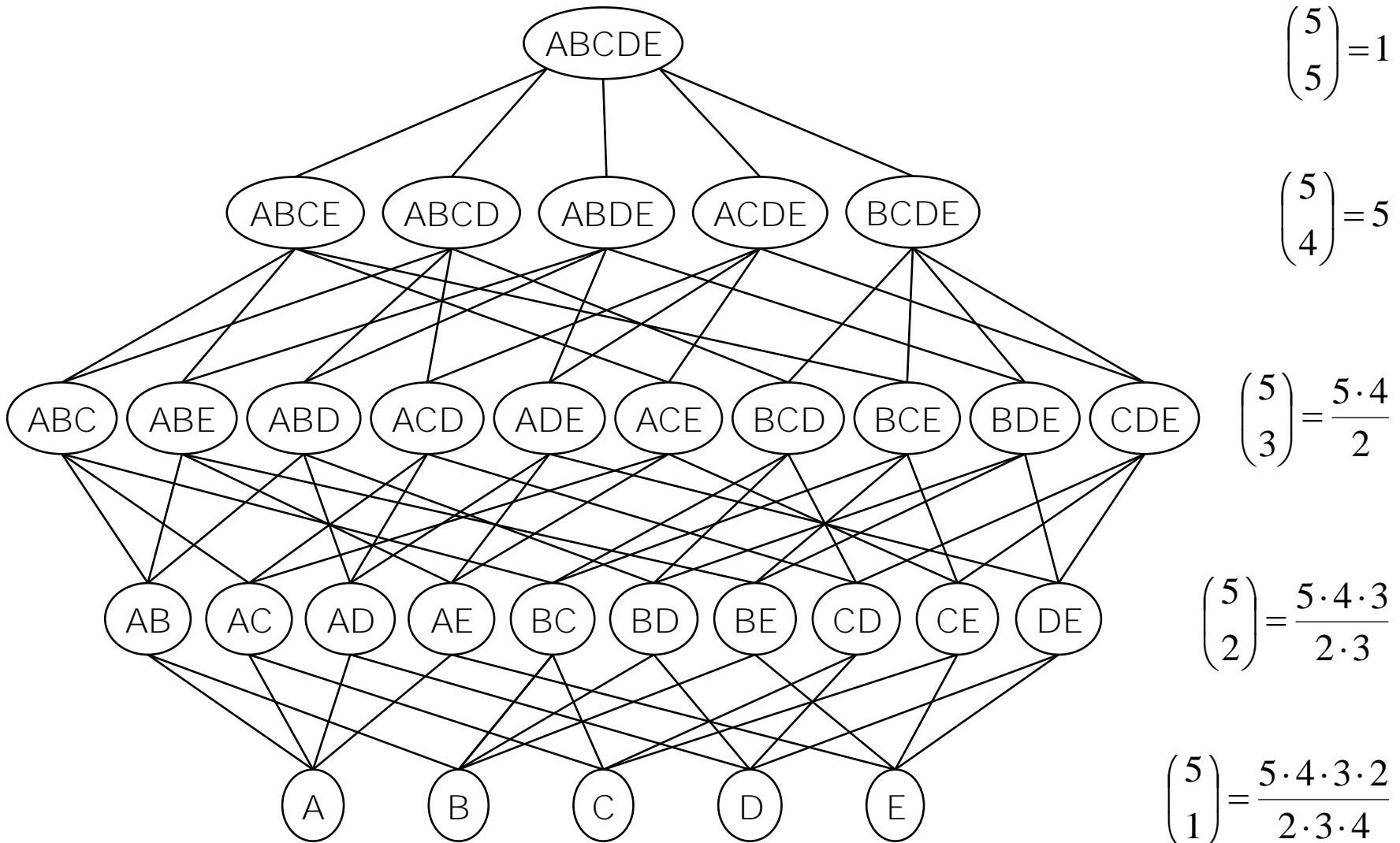


10 columns



Size of the lattice

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Computational feasibility

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For a lattice over n columns

- $\binom{n}{k}$ combinations of size k
- All combinations: $2^n - 1$ (**let's ignore** -1 for the remaining slides)
- Largest solution set: $\binom{n}{\frac{n}{2}}$ minimal uniques are of size $\frac{n}{2}$
 - Verifying minimality, requires to check also all combinations of size $\frac{n}{2} - 1$
- Adding a column doubles search space

Brute forcing Uniprot

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- Data set about proteins with 223 columns
- Combinations: $\sim 1.3 \cdot 10^{67}$
- Largest solution: $\sim 7.2 \cdot 10^{65}$
 - There are roughly 10^{50} atoms on earth
- Assuming all uniques are of size 1-9
 - $$\binom{223}{9} + \binom{223}{8} + \dots \approx 3.3 \cdot 10^{15}$$
 - 1ms verification time results in 100ka processing time

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Null values

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- Null values have a wide range of interpretations
 - Unknown (birth day)
 - Non-applicable (driver license number for kids)
 - Undefined (result of integration/outer join)
- What is the minimal unique for the following data set?

A	B	C	D
a	1	x	1
b	2	y	2
c	3	z	5
d	3	⊥	5
e	⊥	⊥	5

Handling null values #1

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- Depends on the actual application
- To find primary keys
 - Remove all columns with null values
 - Result: { A }

A	B	C	D
a	1	x	1
b	2	y	2
c	3	z	5
d	3	⊥	5
e	⊥	⊥	5

Handling null values #2

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- Depends on the actual application
- To define unique constraints
 - SQL defines grouping for null: $\text{null} \neq \text{null}$
 - Result: {A, C} \rightarrow CD unique
 - A column of nulls is unique!

A	B	C	D
a	1	x	1
b	2	y	2
c	3	z	5
d	3	⊥	5
e	⊥	⊥	5

Handling null values #3

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- Depends on the actual application
- To define unique constraints
 - SQL defines distinctness for null: $\text{null} = \text{null}$
 - Result: {A, BC}

A	B	C	D
a	1	x	1
b	2	y	2
c	3	z	5
d	3	⊥	5
e	⊥	⊥	5

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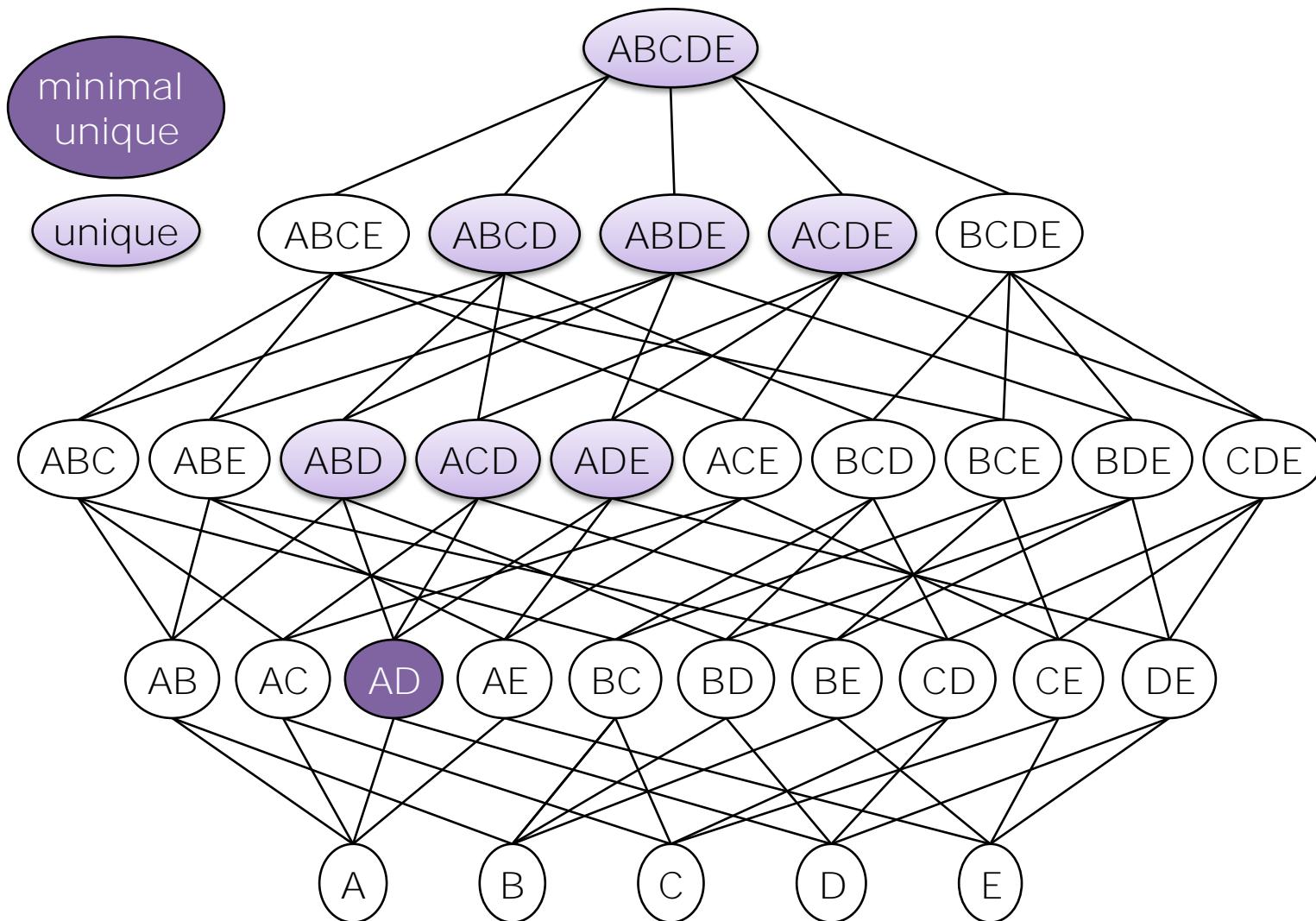
Pruning with uniques

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- Pruning: inferring the type of a combination without actual verification
- If A is unique, supersets must be unique

Pruning effect of a pair

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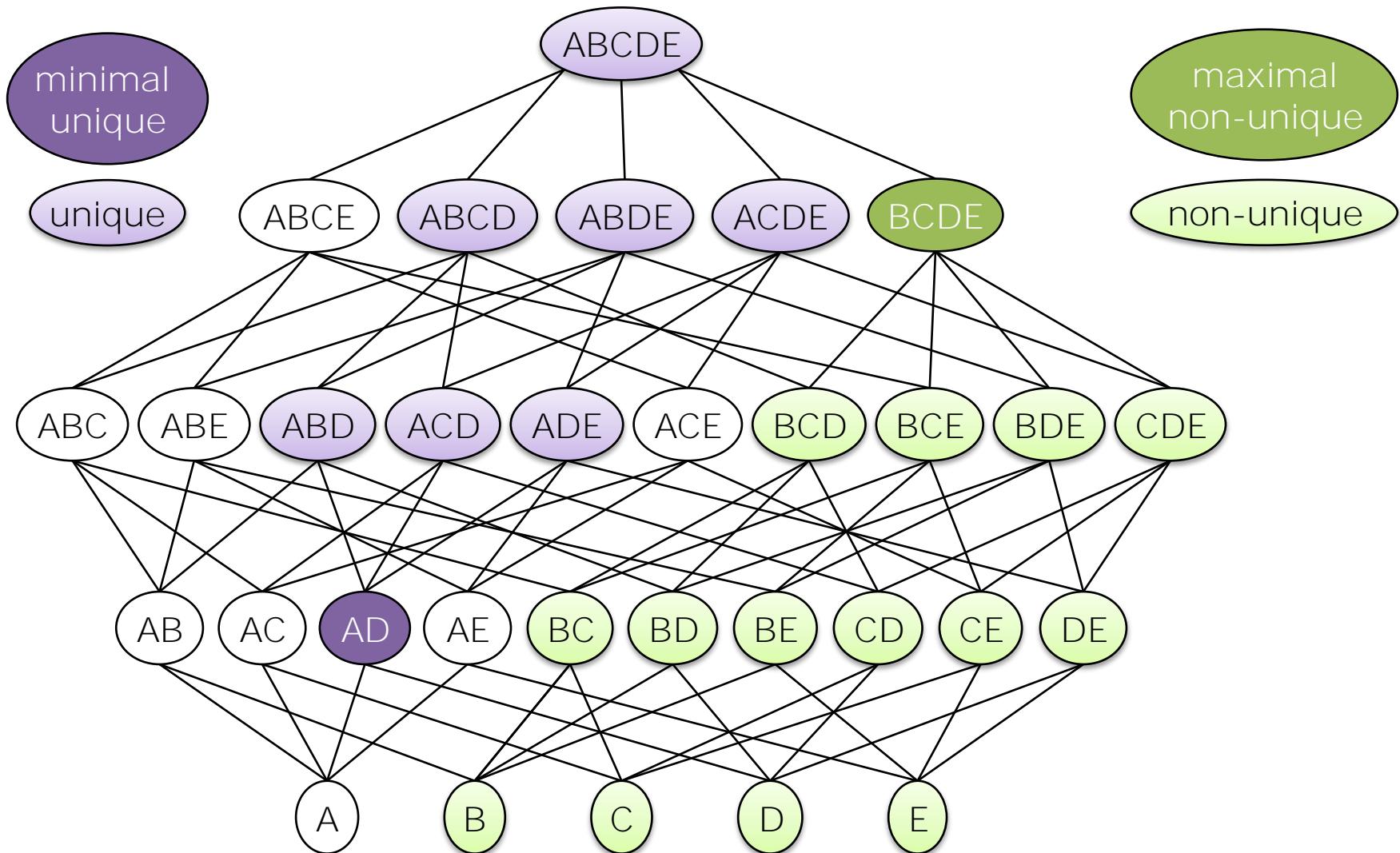
Pruning with uniques #2

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- Pruning: inferring the type of a combination without actual verification
- If A is unique, supersets must be unique
- Finding a unique column prunes half of the lattice
 - Remove column from initial data set and restart
- Finding a unique column pair removes a quarter of the lattice
 - In general, the lattice over the combination is removed
- The pruning power of a combination is reduced by prior findings
 - AB prunes a quarter
 - BC additionally prunes only one eighth
 - ABC already pruned one eights

Pruning both ways

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Pruning on-the-fly

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- Materialization of the lattice is infeasible
 - Only possible for few columns
 - Nodes cannot be removed when discovering unique
- Prune on-the-fly
 - Enumerate nodes as before
 - Skip a node that has been pruned
 - Depending on the approach that might be challenging
 - Might require an efficient index structure
 - Often: candidate generation

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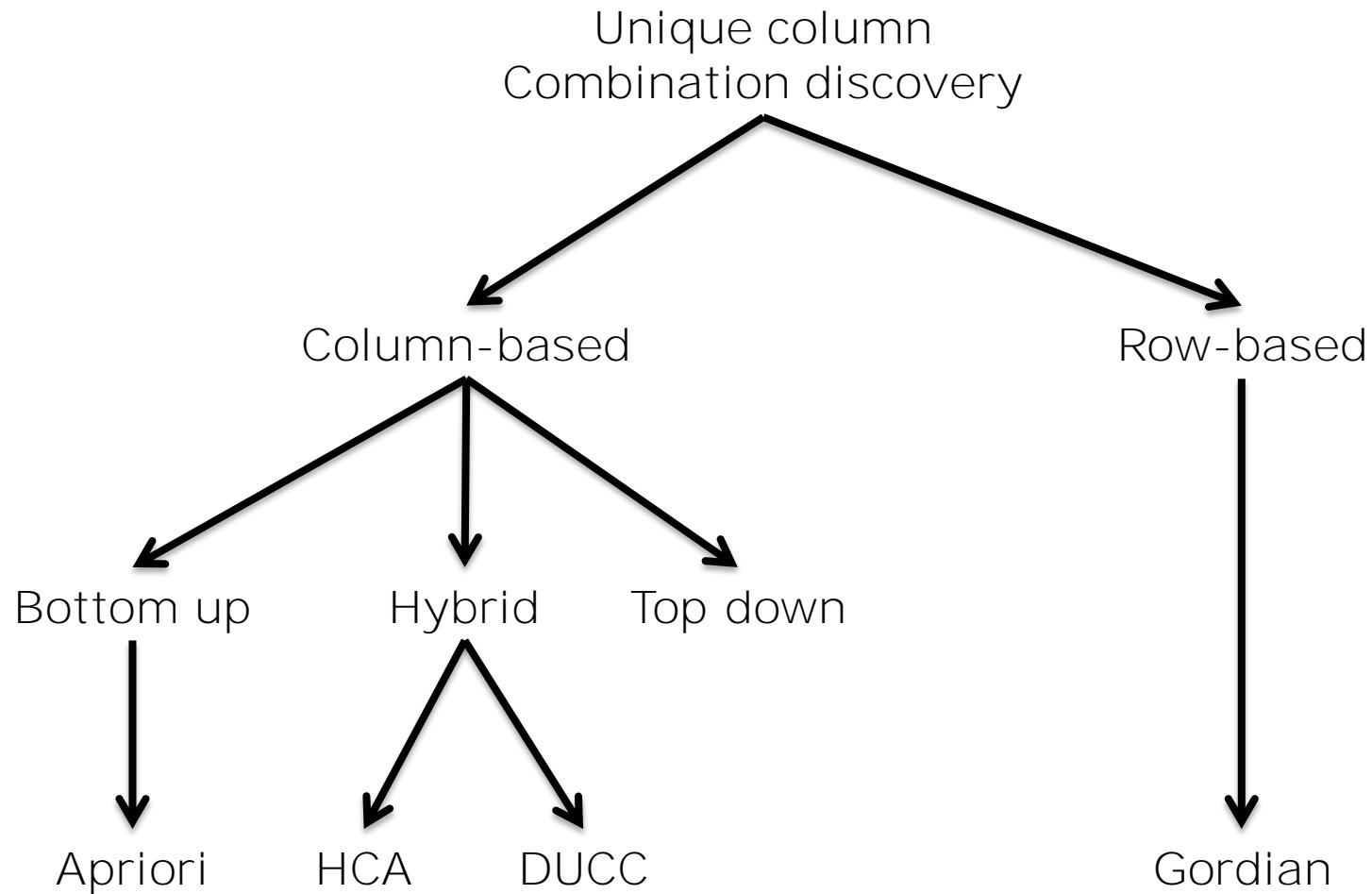
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Discovery Algorithms

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Column-based algorithms

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- Traverse through lattice
- Check for uniqueness
 - Different approaches possible
 - Use database back end and distinctness query
 - ◊ SELECT COUNT(DISTINCT A, B, C) FROM R
 - ◊ Compare with row number
 - Position list indexes (explained later)
 - For now, check is blackbox
- Prune lattice accordingly

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Apriori-based

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- C. Giannella and C. M. Wyss. "Finding minimal keys in a relation instance." (1999).
- Actually does not use much of the apriori idea
- Basic idea:
 - Using the state of combinations of size k
 - We need to visit only unpruned combinations of size $k+1$
- Start with columns
- Check pairs of non-unique columns
- Check triples of non-**unique pairs** ...
- Terminate if no new combinations can be enumerated

Candidate generation

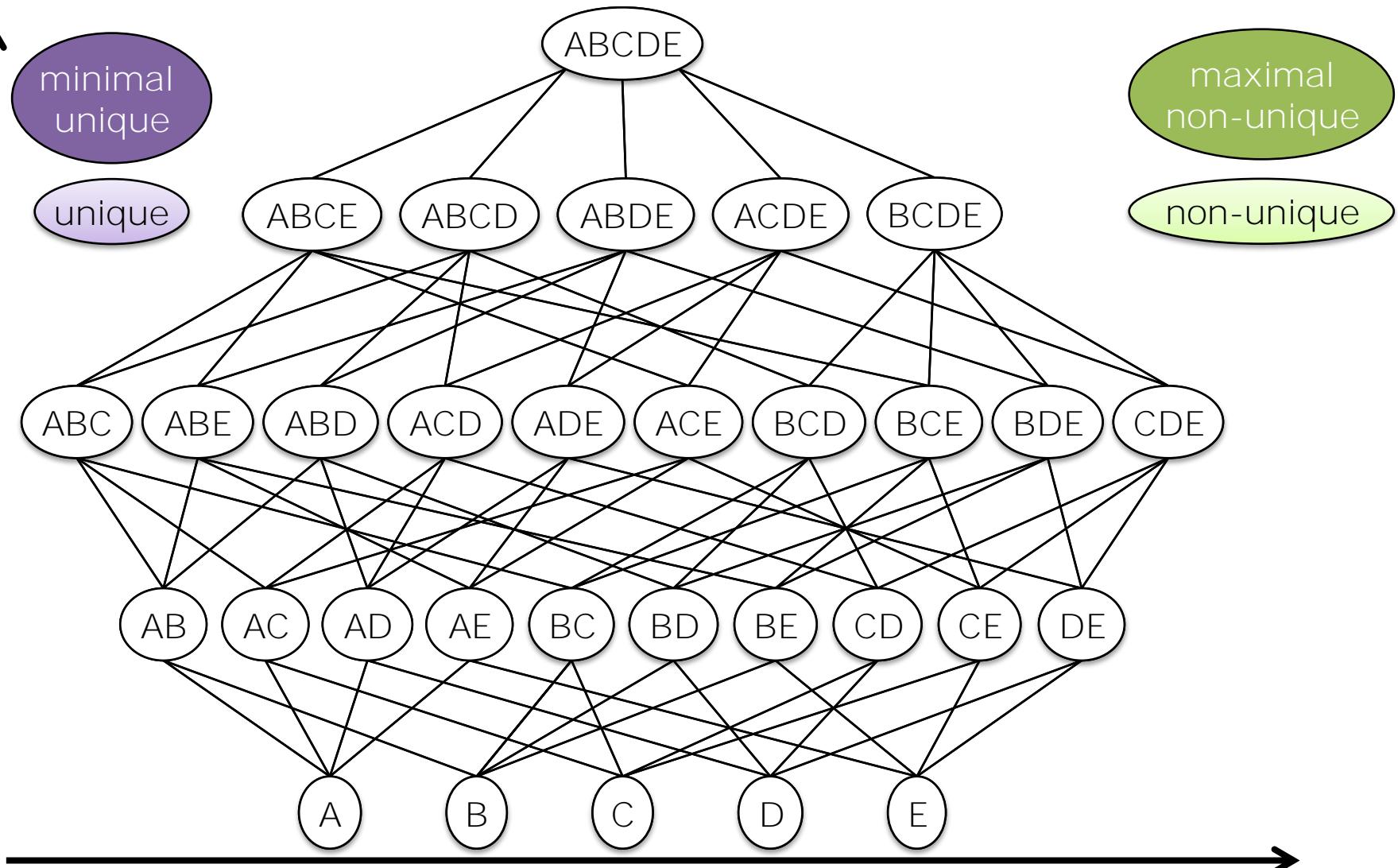
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- Do not generate too many duplicate combinations
- ABC, ABD, ACD, and BCD could point to ABCD

- Apriori: prefix-based generation
- Generate only combination of size n if prefix n-1 matches
- Only ABC and ABD can generate ABCD
 - Still redundant verifications

Apriori visualized

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Characteristics of Apriori

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- Works well for small uniques
 - Bottom-up checks columns first
- Best case: all columns are unique
 - n checks
- Worst case: no uniques = one duplicate row
 - 2^n checks
- Apriori is exponential to n

Extensions

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Top-down

- Start from top and go down
- Performs better if solution set is high up
- Candidate pruning becomes more tricky

Hybrid

- Combine bottom-up and top-down
- Interleave checks
- Works well if solution set has many small and large columns
- Worst case: solution set in the middle

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Histogram-Count-based Apriori

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- Ziawasch Abedjan and Felix Naumann. "Advancing the discovery of unique column combinations." *Proceedings of the international Conference on Information and Knowledge Management*. 2011.
- Extension of bottom-up apriori

- More sophisticated candidate generation
- Uses histograms for pruning
- Finds and uses functional dependencies on-the-fly

HCA candidate generation

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- Maintains a sorted list of non-uniques
 - Avoids duplicate generation of combinations
- Prunes non-minimal uniques efficiently
 - ABC unique, ABD is non-unique
 - ABD would generate ABCD
 - HCA performs quick minimality check with bitsets
- Hybrid approach
 - At least checks if remaining columns contains duplicates

Statistics

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- Prunes column combinations that cannot be unique
 - A and B contains the same value for 4/7 of the data
 - C contains the same value for 5/7 of the data
 - AC cannot be unique, AB might (not very likely)
- Especially viable if there are already indices

A	B	C
1	A	U
1	A	U
1	A	U
1	A	U
2	B	U
3	C	V
4	D	W

Functional dependencies

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- Functional dependency
 - Value of one column determines value of another
 - Birthday->age
- Intuition:
 - If $A \rightarrow B$ and A non-unique, B must be non-unique
 - If $A \rightarrow B$ and B unique, A must also be unique
 - If $A \rightarrow B$ and AC non-unique, BC must be non-unique
 - If $A \rightarrow B$ and BC unique, AC must also be unique
- FD $A \rightarrow B$ can be found with histogram of AB and B
 - Histograms of FDs have the same distinctness counts

Analysis of HCA

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- Works well on data sets with small numbers of columns
- Quickly converges for many small combinations
 - Efficient pruning
- Saves distinctness checks for many pairs

- For larger combinations statistics become less important
 - At some point has to try all combinations
- Suffers from the same general complexity of Apriori

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DUCC

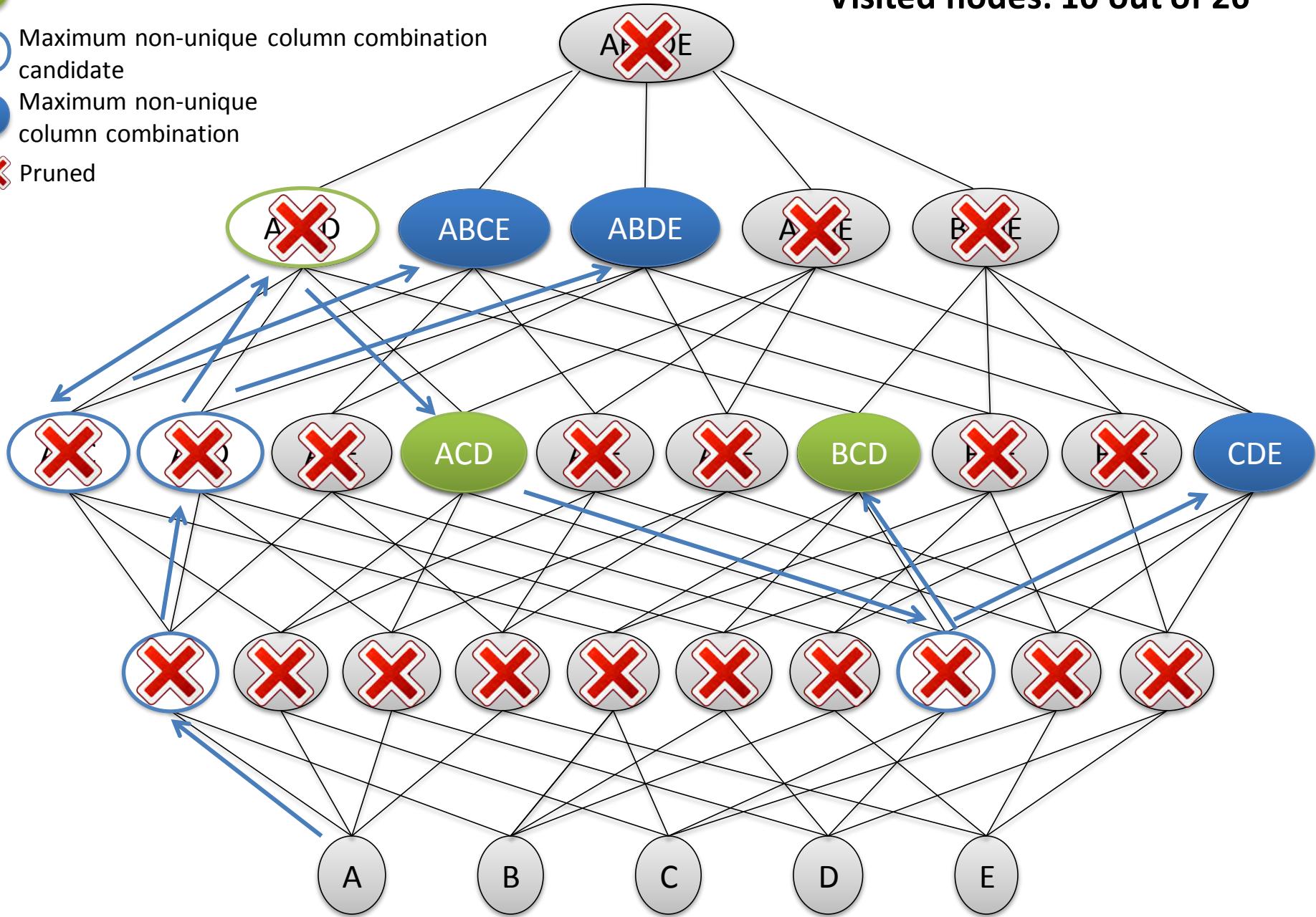
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- Arvid Heise, Jorge-Arnulfo Quiané-Ruiz, Ziawasch Abedjan, Anja Jentzsch, and Felix Naumann, “Scalable Discovery of Unique Column Combinations”, *in preparation*
- Done during internship at QCRI

- Basic idea: random walk through lattice
- Pick random superset if current combination is non-unique
- Pick random subset otherwise
- Lazy prune with previously visited nodes

- Minimum unique column combination candidate
- Minimum unique column combination
- Maximum non-unique column combination candidate
- Maximum non-unique column combination
- ✗ Pruned

Visited nodes: 10 out of 26



Position List Index

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- Incorporates row-based pruning
- Intuition: number of duplicates decrease when going up
 - Many unnecessary rows are checked again and again
- Keep track of duplicates with inverted index
 - A: a->{r₁, r₂, r₃}, b->{r₄, r₅}
 - B: 1->{r₁, r₃}, 2->{r₂, r₅}
- We don't need the actual value
 - A: {{r₁, r₂, r₃}, {r₄, r₅}}
 - B: {{r₁, r₃}, {r₂, r₅}}

A	B
a	1
a	2
a	1
b	3
b	2

PLI Intersection

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Initial PLIs

$$\begin{aligned} A &= \{\{r_1, r_2, r_3\}, \{r_4, r_5\}\} = \{A_1, A_2\} \\ B &= \{\{r_1, r_3\}, \{r_2, r_5\}\} = \{B_1, B_2\} \end{aligned}$$

Build(A)

r_1	A_1
r_2	A_1
r_3	A_1
r_4	A_2
r_5	A_2

Probe(B)

$(A_1, B_1) \rightarrow$	$\{r_1, r_3\}$
$(A_1, B_2) \rightarrow$	$\{r_2\}$
$(A_2, B_2) \rightarrow$	$\{r_5\}$

Consolidate(AB)

$$AB = \{\{r_1, r_3\}\}$$

Analysis of PLI

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- Space complexity: $n \cdot \text{sizeof}(\text{long}) + \frac{n}{2} \cdot \text{sizeof}(\text{array})$
- Intersection time complexity: $O(n+n)$

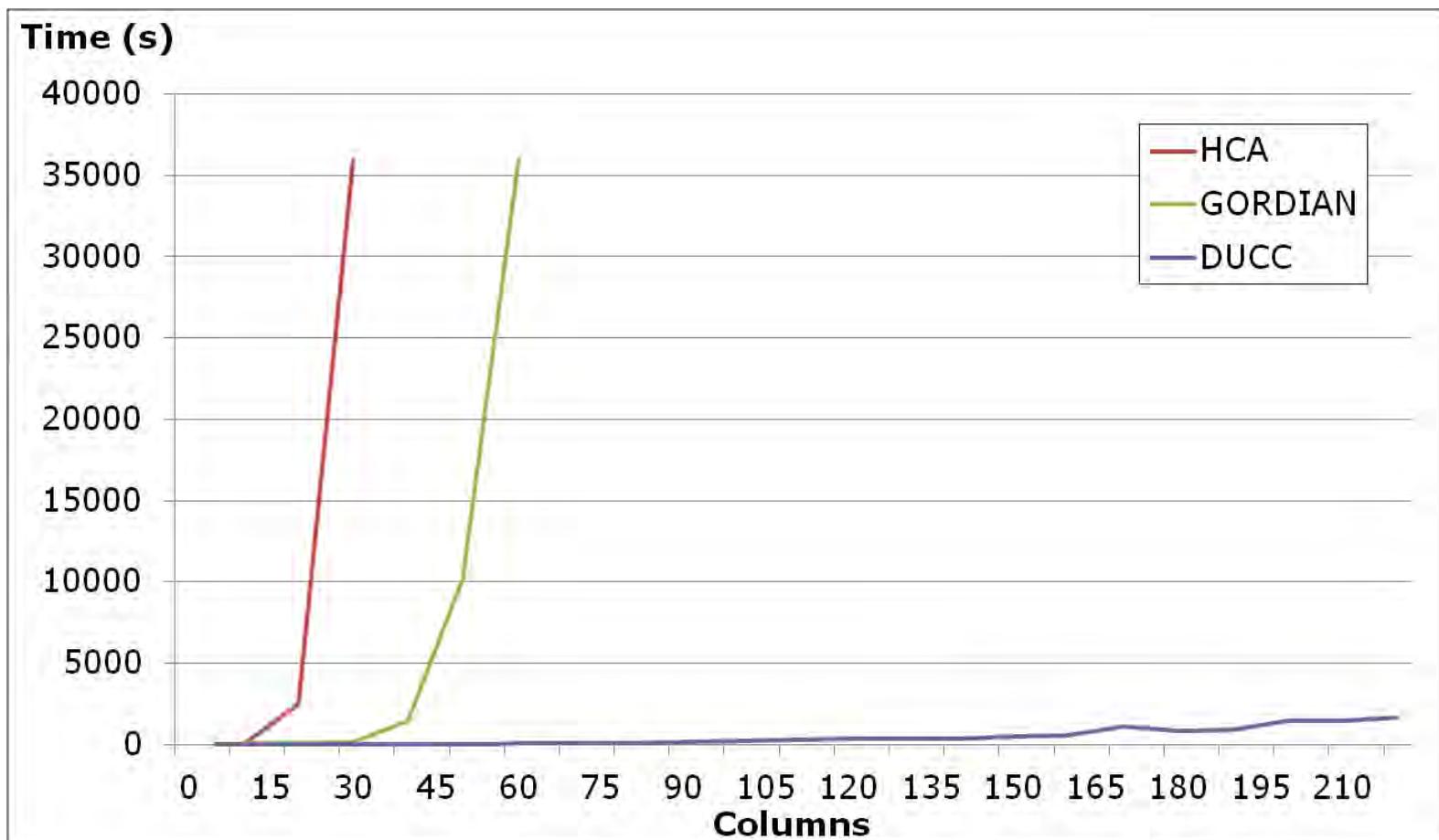
- Hash bigger PLI and probe smaller PLI
- If there is enough main memory
 - Keep PLI of columns in main memory
 - Going up in the lattice requires only to probe the current PLI
 - ◊ Becomes increasingly fast when going up
 - ◊ <1ms for most combinations

- Going down
 - Unfortunately, PLI does not help
 - Start from scratch

Experiments

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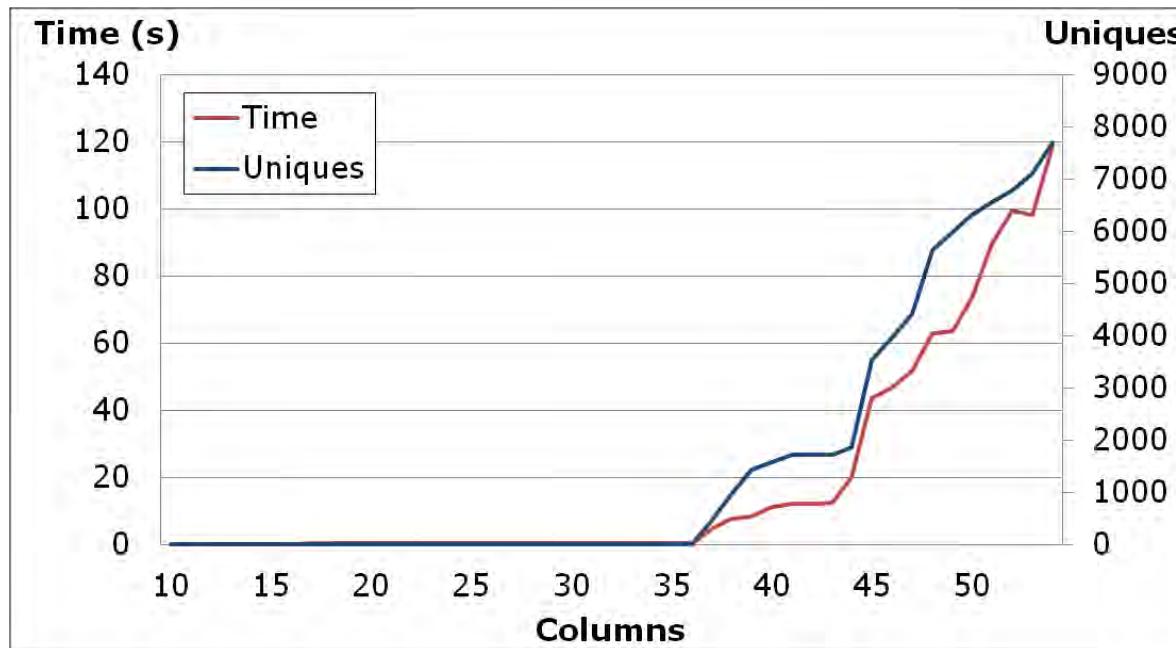
- Uniprot, 100k rows, (DUCC null = null)



Analysis of DUCC

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- Runtime mainly depends on size of solution set



- Worst case: solution set in the middle
- Aggressive pruning may lead to loss of minimal uniques!
 - Gordian's final step can be used to plug these holes

Scaling up and out

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- Scalability is major design goal of DUCC
- Random walk well suited for parallelization
 - Few coordination overhead
- Threads/worker share findings through event bus
 - Uniques/non-uniques
 - Holes in graph
- Lock-free to avoid bottlenecks
 - Only memory barrier in local event bus

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Gordian

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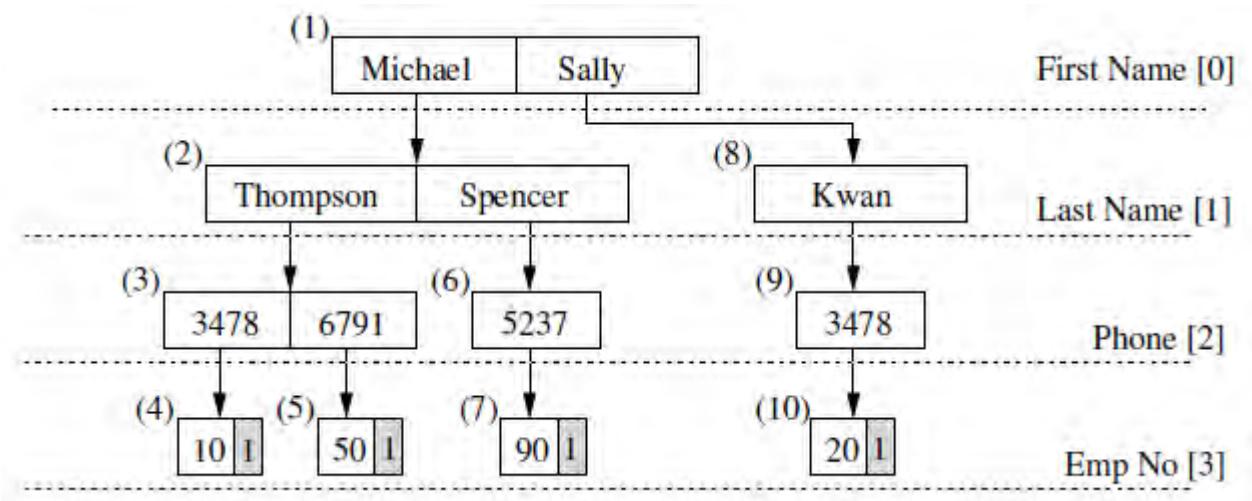
- Yannis Sismanis et al. "GORDIAN: efficient and scalable discovery of composite keys." *Proceedings of the international conference on Very Large Data Bases*. 2006.
- Row-based algorithm

- Builds prefix tree while reading data
- Determines maximal non-uniques
- Compute minimal uniques from maximal non-uniques

Prefix tree

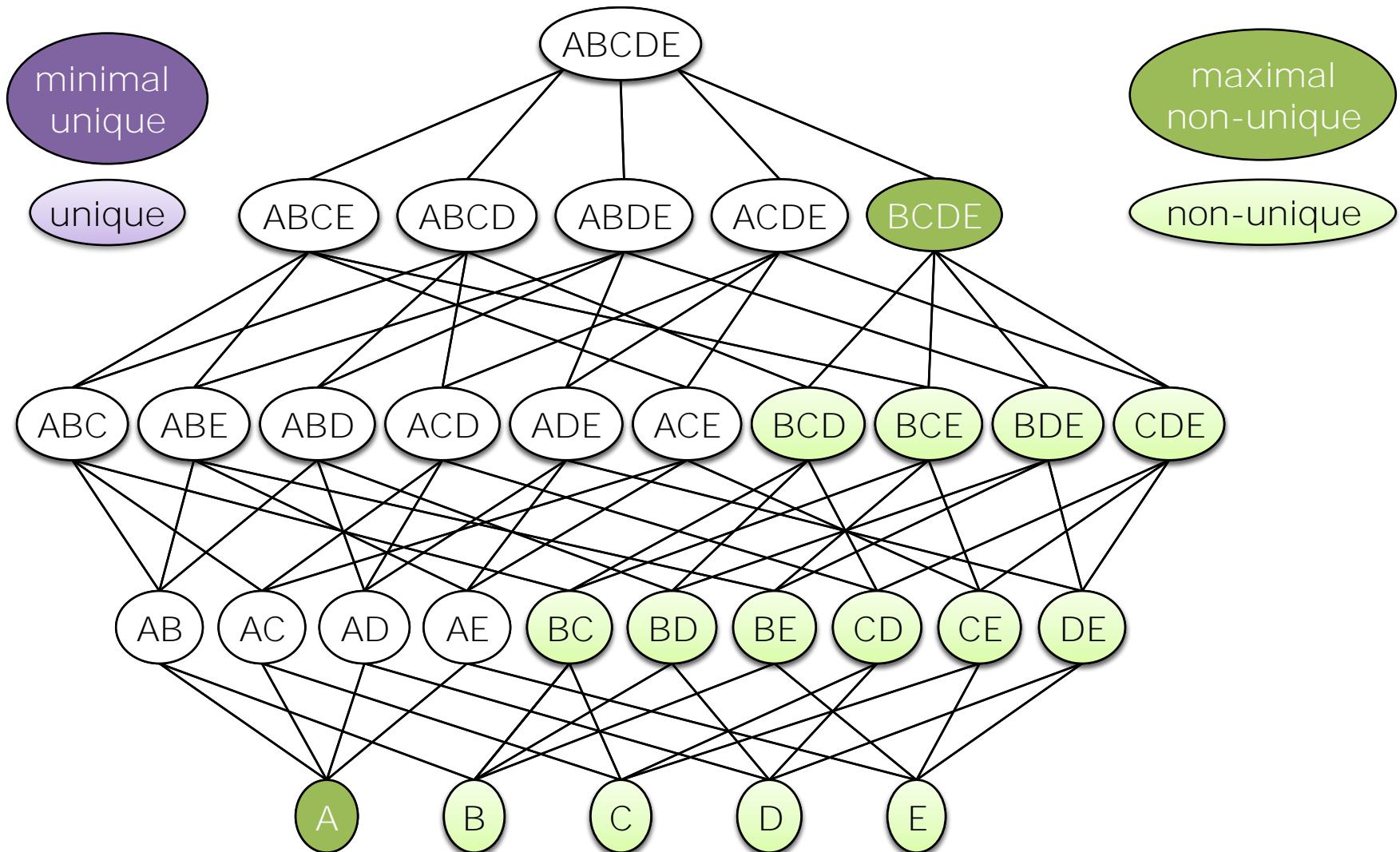
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<i>FirstName</i>	<i>LastName</i>	<i>Phone</i>	<i>EmpNo</i>	<i>COUNT</i>
Michael	Thompson	3478	10	1
Sally	Kwan	3478	20	1
Michael	Spencer	5237	90	1
Michael	Thompson	6791	50	1



Calculating minimal uniques

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Analysis Gordian

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- According to paper, polynomial in the number of tuples for data with a Zipfian distribution of values
 - Can abort scan as soon as duplicate has been found
- Worst case
 - Exponential in the number of columns
 - All data needs to be stored in memory
- Computing minimal uniques from maximal non-uniques
 - $O(\text{uniques}^2 \cdot \text{columns})$
 - Can be sped up with presorted list

Outlook

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- Finding primary keys
 - Uniqueness is necessary criteria
 - No null values
 - Include other features
 - ◆ Name includes “id”, number of columns
- Approximate uniques
 - 99.9% of the data unique
 - Useful to detect data errors
 - Gordian, HCA, and DUCC can be easily modified
- Heuristics with sampling