A fast approach for parallel deduplication on multicore processors

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Overview

- General Blocking
- MD-Approach Overview
- MapReduce Implementation
- Evaluation
- Discussion

General Blocking

DiscID	DiscName	Genre	Year	
1	From The Cradle - Eric Clapton	Blues	1994	
2	Marvin Gaye - Here, My Dear	Soul	1975	
3	The Beatles - A Hard Day's Night	Blues	1964	
4	Eric Clapton - From the Cradle	Blues	1995	
5	Beatles - A Hard Day's Night	Rock	1964	
6	Curtis Mayfield - Curtis	Soul	1970	

General Blocking - Blocking Key

DiscID	DiscName	Genre	Year	
1	From The Cradle - Eric Clapton	Blues	1994	
2	Marvin Gaye - Here, My Dear	Soul	1975	
3	The Beatles - A Hard Day's Night	Blues	1964	
4	Eric Clapton - From the Cradle	Blues	1995	
5	Beatles - A Hard Day's Night	Rock	1964	
6	Curtis Mayfield - Curtis	Soul	1970	

General Blocking - Balance Problem

DiscID	DiscName	Genre	Year	
1	From The Cradle - Eric Clapton	Blues	1994	• •
3	The Beatles - A Hard Day's Night	Blues	1964)
4	Eric Clapton - From the Cradle	Blues	1995	/ /

2	Marvin Gaye - Here, My Dear	Soul	1975	🔪	\langle
6	Curtis Mayfield - Curtis	Soul	1970	🖌	

5	Beatles - A Hard Day's Night	Rock	1964	
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General Blocking - Keys Problem

DiscID	DiscName	Genre	Year	
1	From The Cradle - Eric Clapton	Blues	1994	
2	Marvin Gaye - Here, My Dear	Soul	1975	
3	The Beatles - A Hard Day's Night	Blues	1964	
4	Eric Clapton - From the Cradle	Blues	1995	
5	Beatles - A Hard Day's Night	Rock	1964	
6	Curtis Mayfield - Curtis	Soul	1970	

Blocking Functions & Multipass

- blocking functions are defined as followed:
 - $bf_1(record) = \{genre\}$
 - \circ bf₂(record) = {year, genre}
 - $bf_3(record) = \{1^{st} \ 3 \ letters \ of \ genre, \ 1^{st} \ 3 \ digits \ of \ year\}$
- in a n-multipass several blocking functions are applied to each record

• BFS =
$$\{bf_{1,} bf_{2,} ..., bf_{n}\}$$









MD-Approach - MapReduce Overview



Map-Reduce Implementation Phase I - First Blocking Step

- create dataset segments
- only map phase
- emits key-value pair
 - generated blocking key as key, e.g.
 bf(record) = {1st 3 letters of genre, 1st 3 digits of year}



record as value



Map-Reduce Implementation Phase I - First Blocking Step

• multi-passing

○ set of n several blocking functions

■ BFS =
$$\{bf_{1}, bf_{2}, ..., bf_{n}\}$$

• for each record emit **at once**:

↓ bf ₁ bf	2
bf	
Sou197 2 Marvin Gaye - Here, My Dear Soul 1975	Sou197

MarvSou 2	2	Marvin Gaye - Here, My Dear	Soul	1975		
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Map-Reduce Implementation Phase II - Sort Blocks & Match

- identify unbalanced blocks
 - compare the record count of each block with a threshold
 - use reduce function until a certain threshold is reached



- reduce step (match step)
 - receives all records with the same key (here same block)
 - nested-loop pairwise comparing
 - outputs pairs of similar records

Map-Reduce Implementation Phase III - Second Blocking Step

- only unbalanced blocks
- map: expand blocking key from first blocking step
 - e.g. bf₁(record) = {1st 3 letters of genre, 1st 3 digits of year} → bf₁'(record) = {all letters of genre, all digits of year}
 - creates very fine granular blocks



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	Blu199	4	Eri	ic Clapton - From the Cradl	e E	Blues	199	95)
			i				i				1
	Blues1994		1	From The Cradle - Eric C	Clapton	Blu	es	1994			
X									-		
Blue	es1995	4	Eric Cla	apton - From the Cradle	Blues		1995				

Map-Reduce Implementation Phase III - Second Blocking Step

- to avoid loss of true positives use 'sliding window approach'
 - create an index structure for fine-grained keys after map phase
 - compare with k-nearest neighbors
 - if the similarity is high enough merge records with very similar keys to bigger blocks again
- reduce step (match) is same as in Phase II

Map-Reduce Implementation Phase IV - Merge Pairs

- short map-reduce operations to clean output file
 - identify and remove replicated pairs
 - multipass generates duplicates of detected records



Evaluation

- Phoenix MR framework was used for implementation shared memory-architecture
- synthetic dataset generated by *Febrl* (1M, 2M, 4M, each with 10% duplicates)
- compared with BTO-BK
- used different similarity metrics for different approaches



Relevance for the seminar

- interesting and intuitive main idea
- due to weaknesses in English language, sometimes hard to understand
- the MR-specific implementation details are very rare
- the mapping from a shared-memory (Phoenix) onto a shared-nothing (Hadoop, Stratosphere) architecture will be challenging
- to sum best things up:
 - single-run multi-pass
 - load balancing through re-blocking



1. Dal Bianco, Guilherme, Renata Galante, and Carlos A. Heuser. <u>A fast approach for parallel</u> <u>deduplication on multicore processors</u>. In Proceedings of the ACM Symposium on Applied Computing, 2011.

Map-Reduce Implementation First MR-Step

map-step

emits (blocking-key, value)

- identify unbalanced blocks
- reduce-step (balanced blocks only)
 - similarity function
 - arithmetic average
 - find duplicate by threshold



Map-Reduce Implementation Second MR-Step

- map-step
 - emits expanded blocking-key
- "sliding window sort" (binary search)
- reduce-step
 - same as in First MR-Step

